

Discrete Structuren)
Hertentamen II, 2007

The problems are to be solved within 3 hrs.

The use of supporting material (books, notes, calculators) is not allowed.

In each problem you can obtain 10 points, i.e. 100 in total. Your partial result for the first 5 problems may be replaced by your grade in the midterm exam ($\times 5$), provided the grade was ≥ 5.5 .

Some useful hints:

- Really read these hints.
 - Give precise arguments for all your answers.
 - You can write in English or Dutch, but in any case use a readable font!
 - Counterexamples prove that a statement is not true, but positive examples do not prove general validity.
 - If you refer to the hand-out sheet, numbers of implications etc. are sufficient.
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1. Prove that the following proposition

$$(\neg p \rightarrow r) \leftrightarrow ((r \rightarrow q) \rightarrow (p \vee r))$$

is a tautology. Use the form of an *annotated linear proof* (geannoteerd lineair bewijs).

2. Prove (by cases) that $|x + y| \leq |x| + |y|$ for $x, y \in \mathbb{R}$.

3. Prove by (infinite) mathematical induction: $\sum_{i=0}^n i(i+1) = \frac{n(n+1)(n+2)}{3}$ for $n \in \mathbb{N}$

4. Give an explicit expression for the sequence s_n , defined by

$$\begin{aligned} s_0 &= 1 \\ s_1 &= 1 \\ s_n &= 2s_{n-1} + 2s_{n-2} \quad \text{for } n \geq 2 \end{aligned}$$

5.

- (a) Let $s(n)$ ($n \in \mathbb{N}$) be a sequence. Define the meaning of $s(n) = O(n)$ and of $s(n) = \Theta(n)$.
(b) Are the following statements true or false? (Give precise arguments!)

$$2^{2^n} = O(2^n) \quad 2^{n+1} = \Theta(2^n)$$

6. Let the relation \sim on \mathbb{N} be defined by: $m \sim n$ if and only if $5 \mid (m - n)$ (i.e. 5 divides $m - n$). Show explicitly that \sim satisfies the properties of an equivalence relation. What are the equivalence classes of \sim ?

7.

(a) Show that the proposition

$$[\exists x p(x)] \wedge [\exists x q(x)] \rightarrow \exists x [p(x) \wedge q(x)]$$

is not a tautology. You can do this by giving examples for $p(x)$ and $q(x)$ for which the proposition is false.

(b) Show that the proposition

$$\exists x \forall y p(x, y) \rightarrow \forall x \exists y p(x, y)$$

is not a tautology. Again it is sufficient to show that the proposition is false for a particular $p(x, y)$.

8. Let A be the Boolean matrix

$$A = \begin{bmatrix} 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

(a) Calculate $A * A$.

(b) Is the relation corresponding to A transitive? Explain your answer!

(c) Which matrices represent the symmetric closure, the reflective closure, and the transitive closure of the relation corresponding to A ?

9. After having graduated you have been hired by a manufacturer of computer hardware. Your first task is to specify a scheme for the serial number of a new product. You decide on using alphanumerical characters, i.e. the 26 capital letters and the 10 digits.

Your company does not expect to manufacture more than 1000000000 (i.e. 10^9) of these devices. Out of how many alphanumerical characters should the serial number consist, such that there will be a unique serial number for each manufactured device and the serial number of each device is as short as possible?

You should not give the result as a number; it is sufficient to provide an analytic expression, e.g. $\exp[12]$ instead of 162754.7914....

10.

(a) How many edges are there in a complete graph with $n = 11$ vertices?

(b) How many edges are there in a binary rooted tree with $n = 23$ vertices?

(c) How many edges are there in a ternary rooted tree with $n = 23$ vertices?

1. $\neg\neg p \iff p$	double negation
2a. $(p \vee q) \iff (q \vee p)$ b. $(p \wedge q) \iff (q \wedge p)$ c. $(p \leftrightarrow q) \iff (q \leftrightarrow p)$	commutative laws
3a. $[(p \vee q) \vee r] \iff [p \vee (q \vee r)]$ b. $[(p \wedge q) \wedge r] \iff [p \wedge (q \wedge r)]$	associative laws
4a. $[p \vee (q \wedge r)] \iff [(p \vee q) \wedge (p \vee r)]$ b. $[p \wedge (q \vee r)] \iff [(p \wedge q) \vee (p \wedge r)]$	distributive laws
5a. $(p \vee p) \iff p$ b. $(p \wedge p) \iff p$	idempotent laws
6a. $(p \vee 0) \iff p$ b. $(p \vee 1) \iff 1$ c. $(p \wedge 0) \iff 0$ d. $(p \wedge 1) \iff p$	identity laws ¹
7a. $(p \vee \neg p) \iff 1$ b. $(p \wedge \neg p) \iff 0$	
8a. $\neg(p \vee q) \iff (\neg p \wedge \neg q)$ b. $\neg(p \wedge q) \iff (\neg p \vee \neg q)$ c. $(p \vee q) \iff \neg(\neg p \wedge \neg q)$ d. $(p \wedge q) \iff \neg(\neg p \vee \neg q)$	DeMorgan laws
9. $(p \rightarrow q) \iff (\neg q \rightarrow \neg p)$	contrapositive
10a. $(p \rightarrow q) \iff (\neg p \vee q)$ b. $(p \rightarrow q) \iff \neg(p \wedge \neg q)$	implication
11a. $(p \vee q) \iff (\neg p \rightarrow q)$ b. $(p \wedge q) \iff \neg(p \rightarrow \neg q)$	
12a. $[(p \rightarrow r) \wedge (q \rightarrow r)] \iff [(p \vee q) \rightarrow r]$ b. $[(p \rightarrow q) \wedge (p \rightarrow r)] \iff [p \rightarrow (q \wedge r)]$	
13. $(p \leftrightarrow q) \iff [(p \rightarrow q) \wedge (q \rightarrow p)]$	equivalence
14. $[(p \wedge q) \rightarrow r] \iff [p \rightarrow (q \rightarrow r)]$	exportation law
15. $(p \rightarrow q) \iff [(p \wedge \neg q) \rightarrow 0]$	reductio ad absurdum

16. $p \implies (p \vee q)$	addition
17. $(p \wedge q) \implies p$	simplification
18. $(p \rightarrow 0) \implies \neg p$	absurdity
19. $[p \wedge (p \rightarrow q)] \implies q$	modus ponens
20. $[(p \rightarrow q) \wedge \neg q] \implies \neg p$	modus tollens
21. $[(p \vee q) \wedge \neg p] \implies q$	disjunctive syllogism
22. $p \implies [q \rightarrow (p \wedge q)]$	
23. $[(p \leftrightarrow q) \wedge (q \leftrightarrow r)] \implies (p \leftrightarrow r)$	transitivity of \leftrightarrow
24. $[(p \rightarrow q) \wedge (q \rightarrow r)] \implies (p \rightarrow r)$	transitivity of \rightarrow or hypothetical syllogism
25a. $(p \rightarrow q) \implies [(p \vee r) \rightarrow (q \vee r)]$ b. $(p \rightarrow q) \implies [(p \wedge r) \rightarrow (q \wedge r)]$ c. $(p \rightarrow q) \implies [(q \rightarrow r) \rightarrow (p \rightarrow r)]$	
26a. $[(p \rightarrow q) \wedge (r \rightarrow s)] \implies [(p \vee r) \rightarrow (q \vee s)]$ b. $[(p \rightarrow q) \wedge (r \rightarrow s)] \implies [(p \wedge r) \rightarrow (q \wedge s)]$	constructive dilemmas
27a. $[(p \rightarrow q) \wedge (r \rightarrow s)] \implies [(\neg q \vee \neg s) \rightarrow (\neg p \vee \neg r)]$ b. $[(p \rightarrow q) \wedge (r \rightarrow s)] \implies [(\neg q \wedge \neg s) \rightarrow (\neg p \wedge \neg r)]$	destructive dilemmas

28a. $\neg 1 \iff 0$	0-1-wetten
b. $\neg 0 \iff 1$	
29a. $(p \wedge (p \vee q)) \iff p$	absorptie
b. $(p \vee (p \wedge q)) \iff p$	

Equivalenties:

30a.	$p \iff \forall x p$ (x niet vrij in p)	loze kwantificatie
ba.	$p \iff \exists x p$ (x niet vrij in p)	
31a.	$\forall x p(x) \iff \forall y p(y)$	herbenoemen van gebonden variabele
b.	$\exists x p(x) \iff \exists y p(y)$	
32a.	$\forall x \forall y p(x, y) \iff \forall y \forall x p(x, y)$	kwantorwisseling
b.	$\exists x \exists y p(x, y) \iff \exists y \exists x p(x, y)$	
33a.	$\forall x(p(x) \wedge q(x)) \iff \forall x p(x) \wedge \forall x q(x)$	$\forall(V) \iff \forall A \wedge A$ (A distribueert over \forall)
b.	$\exists x(p(x) \vee q(x)) \iff \exists x p(x) \vee \exists x q(x)$	$\exists(V) \iff \exists E \vee E$ (E distribueert over \vee)
34a.	$\forall x(p \vee q(x)) \iff p \vee \forall x q(x)$ (x niet vrij in p)	
b.	$\exists x(p \wedge q(x)) \iff p \wedge \exists x q(x)$ (x niet vrij in p)	
35a.	$\neg \forall x p(x) \iff \exists x \neg p(x)$	$\neg \forall \iff \exists \neg$ (De Morgan)
b.	$\neg \exists x p(x) \iff \forall x \neg p(x)$	$\neg \exists \iff \forall \neg$ (De Morgan)

Implicaties:

36.	$\forall x p(x) \iff \exists x p(x)$	$A \iff E$
37.	$\exists x \forall y p(x, y) \iff \forall y \exists x p(x, y)$	$EA \iff AE$
38a.	$\forall x p(x) \vee \forall x q(x) \iff \forall x(p(x) \vee q(x))$	$\forall A \vee A \iff \forall(A \vee)$
b.	$\exists x p(x) \vee \exists x q(x) \iff \exists x(p(x) \vee q(x))$	$\exists E \vee E \iff \exists(E \vee)$